

Non-degenerate Hilbert Cubes in Random Sets

par CSABA SÁNDOR

RÉSUMÉ. Une légère modification de la démonstration du lemme des cubes de Szemerédi donne le résultat plus précis suivant: si une partie S de $\{1, \dots, n\}$ vérifie $|S| \geq \frac{n}{2}$, alors S contient un cube de Hilbert non dégénéré de dimension $\lfloor \log_2 \log_2 n - 3 \rfloor$. Dans cet article nous montrons que dans un ensemble aléatoire avec les probabilités $\Pr\{s \in S\} = 1/2$ indépendantes pour $1 \leq s \leq n$, la plus grande dimension d'un cube de Hilbert non dégénéré est *proche* de $\log_2 \log_2 n + \log_2 \log_2 \log_2 n$ presque sûrement et nous déterminons la fonction seuil pour avoir un k -cube non dégénéré.

ABSTRACT. A slight modification of the proof of Szemerédi's cube lemma gives that if a set $S \subset [1, n]$ satisfies $|S| \geq \frac{n}{2}$, then S must contain a non-degenerate Hilbert cube of dimension $\lfloor \log_2 \log_2 n - 3 \rfloor$. In this paper we prove that in a random set S determined by $\Pr\{s \in S\} = \frac{1}{2}$ for $1 \leq s \leq n$, the maximal dimension of non-degenerate Hilbert cubes is a.e. nearly $\log_2 \log_2 n + \log_2 \log_2 \log_2 n$ and determine the threshold function for a non-degenerate k -cube.

1. Introduction

Throughout this paper we use the following notations: let $[1, n]$ denote the first n positive integers. The coordinates of the vector $\mathbf{A}^{(k,n)} = (a_0, a_1, \dots, a_k)$ are selected from the positive integers such that $\sum_{i=0}^k a_i \leq n$. The vectors $\mathbf{B}^{(k,n)}$, $\mathbf{A}_i^{(k,n)}$ are interpreted similarly. The set S_n is a subset of $[1, n]$. The notations $f(n) = o(g(n))$ means $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$. An arithmetic progression of length k is denoted by AP_k . The rank of a matrix A over the field \mathbb{F} is denoted by $r_{\mathbb{F}}(A)$. Let \mathbb{R} denote the set of real numbers and \mathbb{F}_2 for the finite field of order 2.

Let n be a positive integer, $0 \leq p_n \leq 1$. The random set $S(n, p_n)$ is the random variable taking its values in the set of subsets of $[1, n]$ with the law determined by the independence of the events $\{k \in S(n, p_n)\}$, $1 \leq k \leq n$ with the probability $\Pr\{k \in S(n, p_n)\} = p_n$. This model is often used for proving the existence of certain sequences. Given any combinatorial number theoretic property P , there is a probability that $S(n, p_n)$ satisfies P , which we write $\Pr\{S(n, p_n) \models P\}$. The function $r(n)$ is called a threshold function for a combinatorial number theoretic property P if

- (i) When $p_n = o(r(n))$, $\lim_{n \rightarrow \infty} \Pr\{S(n, p_n) \models P\} = 0$,

(ii) When $r(n) = o(p(n))$, $\lim_{n \rightarrow \infty} \Pr\{S(n, p_n) \models P\} = 1$,

or visa versa. It is clear that threshold functions are not unique. However, threshold functions are unique within factors $m(n)$, $0 < \liminf_{n \rightarrow \infty} m(n) \leq \limsup_{n \rightarrow \infty} m(n) < \infty$, that is if p_n is a threshold function for P then p'_n is also a threshold function iff $p_n = O(p'_n)$ and $p'_n = O(p_n)$. In this sense we can speak of the threshold function of a property.

We call $H \subset [1, n]$ a Hilbert cube of dimension k or, simply, a k -cube if there is a vector $\mathbf{A}^{(k, n)}$ such that

$$H = \mathbf{H}_{\mathbf{A}^{(k, n)}} = \{a_0 + \sum_{i=1}^k \epsilon_i a_i : \epsilon_i \in \{0, 1\}\}.$$

The positive integers a_1, \dots, a_k are called the generating elements of the Hilbert cube. The k -cube is non-degenerate if $|H| = 2^k$ i.e. the vertices of the cube are distinct, otherwise it is called degenerate. The maximal dimension of a non-degenerate Hilbert cube in S_n is denoted by $H_{max}(S_n)$, i.e. $H_{max}(S_n)$ is the largest integer l such that there exists a vector $\mathbf{A}^{(l, n)}$ for which the non-degenerate Hilbert cube $\mathbf{H}_{\mathbf{A}^{(l, n)}} \subset S_n$.

Hilbert originally proved that if the positive integers are colored with finitely many colors then one color class contains a k -cube. The density version of theorem was proved by Szemerédi and has since become known as "Szemerédi's cube lemma". The best known result is due to Gunderson and Rödl (see [3]):

Theorem 1.1 (Szemerédi). *For every $d \geq 3$ there exists $n_0 \leq (2^d - 2/\ln 2)^2$ so that, for every $n \geq n_0$, if $A \subset [1, n]$ satisfies $|A| \geq 2n^{1-\frac{1}{2^{d-1}}}$, then A contains a d -cube.*

A direct consequence is the following:

Corollary 1.2. *Every subset S_n such that $|S_n| \geq \frac{n}{2}$ contains a $\lfloor \log_2 \log_2 n \rfloor$ -cube.*

A slight modification of the proof gives that the above set S_n must contain a non-degenerate $\lfloor \log_2 \log_2 n - 3 \rfloor$ -cube.

Obviously, a sequence S has the Sidon property (that is the sums $s_i + s_j$, $s_i \leq s_j$, $s_i, s_j \in S$ are distinct) iff S contains no 2-cube. Godbole, Janson, Locantore and Rapoport studied the threshold function for the Sidon property and gave the exact probability distribution in 1999 (see [2]):

Theorem 1.3 (Godbole, Janson, Locantore and Rapoport). *Let $c > 0$ be arbitrary. Let P denote the Sidon property. Then with $p_n = cn^{-3/4}$,*

$$\lim_{n \rightarrow \infty} \Pr\{S(n, p_n) \models P\} = e^{-\frac{c^4}{12}}.$$

Clearly, a subset $H \subset [1, n]$ is a degenerate 2-cube iff it is an AP_3 . Moreover, an easy argument gives that the threshold function for the event " AP_3 -free" is $p_n = n^{-2/3}$. Hence

Corollary 1.4. *Let $c > 0$ be arbitrary. Then with $p_n = cn^{-3/4}$,*

$$\lim_{n \rightarrow \infty} \Pr\{S(n, p_n) \text{ contains no non-degenerate 2-cube}\} = e^{-\frac{c^4}{12}}.$$

In Theorem 1.5 we extend the previous Corollary.

Theorem 1.5. *For any real number $c > 0$ and any integer $k \geq 2$, if $p_n = cn^{-\frac{k+1}{2^k}}$ then*

$$\lim_{n \rightarrow \infty} \Pr\{S(n, p_n) \text{ contains no non-degenerate } k\text{-cube}\} = e^{-\frac{c^k}{(k+1)!k!}}.$$

In the following we shall find bounds on the maximal dimension of non-degenerate Hilbert cubes in the random subset $S(n, \frac{1}{2})$. Let

$$D_n(\epsilon) = \lfloor \log_2 \log_2 n + \log_2 \log_2 \log_2 n + \frac{(1 - \epsilon) \log_2 \log_2 \log_2 n}{\log 2 \log_2 \log_2 n} \rfloor$$

and

$$E_n(\epsilon) = \lfloor \log_2 \log_2 n + \log_2 \log_2 \log_2 n + \frac{(1 + \epsilon) \log_2 \log_2 \log_2 n}{\log 2 \log_2 \log_2 n} \rfloor.$$

The next theorem implies that for almost all n , $H_{\max}(S(n, \frac{1}{2}))$ concentrates on a single value because for every $\epsilon > 0$, $D_n(\epsilon) = E_n(\epsilon)$ except for a sequence of zero density.

Theorem 1.6. *For every $\epsilon > 0$*

$$\lim_{n \rightarrow \infty} \Pr\{D_n(\epsilon) \leq H_{\max}(S(n, \frac{1}{2})) \leq E_n(\epsilon)\} = 1.$$

2. Proofs

In order to prove the theorems we need some lemmas.

Lemma 2.1. *For $k_n = o(\frac{\log n}{\log \log n})$ the number of non-degenerate k_n -cubes in $[1, n]$ is $(1 + o(1)) \binom{n}{k_n+1} \frac{1}{k_n!}$, as $n \rightarrow \infty$.*

Proof. All vectors $\mathbf{A}^{(k_n, n)}$ are in 1-1 correspondence with all vectors $(v_0, v_1, \dots, v_{k_n})$ with $1 \leq v_1 < v_2 < \dots < v_{k_n} \leq n$ in \mathbb{R}^{k_n+1} according to the formulas $(a_0, a_1, \dots, a_{k_n}) \mapsto (v_0, v_1, \dots, v_{k_n}) = (a_0, a_0 + a_1, \dots, a_0 + a_1 + \dots + a_{k_n})$; and $(v_0, v_1, \dots, v_{k_n}) \mapsto (a_0, a_1, \dots, a_{k_n}) = (v_0, v_1 - v_0, \dots, v_{k_n} - v_{k_n-1})$. Consequently,

$$\binom{n}{k_n+1} = |\{\mathbf{A}^{(k_n, n)} : \mathbf{H}_{\mathbf{A}^{(k_n, n)}} \text{ is non-degenerate}\}| + |\{\mathbf{A}^{(k_n, n)} : \mathbf{H}_{\mathbf{A}^{(k_n, n)}} \text{ is degenerate}\}|.$$

By the definition of a non-degenerate cube we have

$$|\{\mathbf{A}^{(k_n, n)} : \mathbf{H}_{\mathbf{A}^{(k_n, n)}} \text{ is non-degenerate}\}| = k_n! |\{\text{non-degenerate } k_n\text{-cubes in } [1, n]\}|,$$

because permutations of a_1, \dots, a_k give the same k_n -cube. It remains to verify that the number of vectors $\mathbf{A}^{(k_n, n)}$ which generate degenerate k_n -cubes is $o(\binom{n}{k_n+1})$. Let $\mathbf{A}^{(k_n, n)}$ be a vector for which $\mathbf{H}_{\mathbf{A}^{(k_n, n)}}$ is a degenerate k_n -cube. Then there exist integers $1 \leq u_1 < u_2 < \dots < u_s \leq k_n$, $1 \leq v_1 < v_2 < \dots < v_t \leq k_n$ such that

$$a_0 + a_{u_1} + \dots + a_{u_s} = a_0 + a_{v_1} + \dots + a_{v_t},$$

where we may assume that the indices are distinct, therefore $s + t \leq k_n$. Then the equation

$$x_1 + x_2 + \dots + x_s - x_{s+1} - \dots - x_{s+t} = 0$$

can be solved over the set $\{a_1, a_2, \dots, a_{k_n}\}$. The above equation has at most $n^{s+t-1} \leq n^{k_n-1}$ solutions over $[1, n]$. Since we have at most k_n^2 possibilities for (s, t) and at most n possibilities for a_0 , therefore the number of vectors $\mathbf{A}^{(k_n, n)}$ for which $\mathbf{H}_{\mathbf{A}^{(k_n, n)}}$ is degenerate is at most $k_n^2 n^{k_n} = o\left(\binom{n}{k_n+1}\right)$. \square

In the remaining part of this section the Hilbert cubes are non-degenerate.

The proofs of Theorem 1.5 and 1.6 will be based on the following definition. For two intersecting k -cubes $\mathbf{H}_{\mathbf{A}^{(k, n)}}$, $\mathbf{H}_{\mathbf{B}^{(k, n)}}$ let $\mathbf{H}_{\mathbf{A}^{(k, n)}} \cap \mathbf{H}_{\mathbf{B}^{(k, n)}} = \{c_1, \dots, c_m\}$ with $c_1 < \dots < c_m$, where

$$c_d = a_0 + \sum_{l=1}^k \alpha_{d,l} a_l = b_0 + \sum_{l=1}^k \beta_{d,l} b_l, \quad \alpha_{d,l}, \beta_{d,l} \in \{0, 1\} \quad \text{for } 1 \leq d \leq m \text{ and } 1 \leq l \leq k.$$

The rank of the intersection of two k -cubes $\mathbf{H}_{\mathbf{A}^{(k, n)}}$, $\mathbf{H}_{\mathbf{B}^{(k, n)}}$ is defined as follows: we say that $r(\mathbf{H}_{\mathbf{A}^{(k, n)}}, \mathbf{H}_{\mathbf{B}^{(k, n)}}) = (s, t)$ if for the matrices $A = (\alpha_{d,l})_{m \times k}$, $B = (\beta_{d,l})_{m \times k}$ we have $r_{\mathbb{R}}(A) = s$ and $r_{\mathbb{R}}(B) = t$. The matrices A and B are called matrices of the common vertices of $\mathbf{H}_{\mathbf{A}^{(k, n)}}$, $\mathbf{H}_{\mathbf{B}^{(k, n)}}$.

Lemma 2.2. *The condition $r(\mathbf{H}_{\mathbf{A}^{(k, n)}}, \mathbf{H}_{\mathbf{B}^{(k, n)}}) = (s, t)$ implies that $|\mathbf{H}_{\mathbf{A}^{(k, n)}} \cap \mathbf{H}_{\mathbf{B}^{(k, n)}}| \leq 2^{\min\{s, t\}}$.*

Proof. We may assume that $s \leq t$. The inequality $|\mathbf{H}_{\mathbf{A}^{(k, n)}} \cap \mathbf{H}_{\mathbf{B}^{(k, n)}}| \leq 2^s$ is obviously true for $s = k$. Let us suppose that $s < k$ and the number of common vertices is greater than 2^s . Then the corresponding $(0-1)$ -matrices A and B have more than 2^s different rows, therefore $r_{\mathbb{F}_2}(A) > s$, but we know from elementary linear algebra that for an arbitrary $(0-1)$ -matrix M we have $r_{\mathbb{F}_2}(M) \geq r_{\mathbb{R}}(M)$, which is a contradiction. \square

Lemma 2.3. *Let us suppose that the sequences $\mathbf{A}^{(k, n)}$ and $\mathbf{B}^{(k, n)}$ generate non-degenerate k -cubes. Then*

$$(1) \left| \{(\mathbf{A}^{(k, n)}, \mathbf{B}^{(k, n)}) : r(\mathbf{H}_{\mathbf{A}^{(k, n)}}, \mathbf{H}_{\mathbf{B}^{(k, n)}}) = (s, t)\} \right| \leq 2^{2k^2} \binom{n}{k+1} n^{k+1-\max\{s, t\}}$$

for all $0 \leq s, t \leq k$;

$$(2) \left| \{(\mathbf{A}^{(k, n)}, \mathbf{B}^{(k, n)}) : r(\mathbf{H}_{\mathbf{A}^{(k, n)}}, \mathbf{H}_{\mathbf{B}^{(k, n)}}) = (r, r), |\mathbf{H}_{\mathbf{A}^{(k, n)}} \cap \mathbf{H}_{\mathbf{B}^{(k, n)}}| = 2^r\} \right| \leq 2^{2k^2} \binom{n}{k+1} n^{k-r}$$

for all $0 \leq r < k$;

$$(3) \left| \{(\mathbf{A}^{(k, n)}, \mathbf{B}^{(k, n)}) : r(\mathbf{H}_{\mathbf{A}^{(k, n)}}, \mathbf{H}_{\mathbf{B}^{(k, n)}}) = (k, k), |\mathbf{H}_{\mathbf{A}^{(k, n)}} \cap \mathbf{H}_{\mathbf{B}^{(k, n)}}| > 2^{k-1}\} \right| \leq 2^{2k^2+2k} \binom{n}{k+1}.$$

Proof. (1): We may assume that $s \leq t$. In this case we have to prove that the number of corresponding pairs $(\mathbf{A}^{(k, n)}, \mathbf{B}^{(k, n)})$ is at most $\binom{n}{k+1} 2^{2k^2} n^{k+1-t}$. We have already seen in the proof of Lemma 1 that the number of vectors $\mathbf{A}^{(k, n)}$ is at most $\binom{n}{k+1}$. Fix a vector $\mathbf{A}^{(k, n)}$ and count the suitable vectors $\mathbf{B}^{(k, n)}$. Then the matrix B has t linearly independent rows, namely $r_{\mathbb{R}}((\beta_{d_i, l})_{t \times k}) = t$, for some $1 \leq d_1 < \dots < d_t \leq m$, where

$$a_0 + \sum_{l=1}^k \alpha_{d_i, l} a_l = b_0 + \sum_{l=1}^k \beta_{d_i, l} b_l, \quad \alpha_{d_i, l}, \beta_{d_i, l} \in \{0, 1\} \quad \text{for } 1 \leq i \leq t.$$

The number of possible b_0 s is at most n . For fixed $b_0, \alpha_{d_i, l}, \beta_{d_i, l}$ let us study the system of equations

$$a_0 + \sum_{l=1}^k \alpha_{d_i, l} a_l = b_0 + \sum_{l=1}^k \beta_{d_i, l} x_l, \quad \alpha_{d_i, l}, \beta_{d_i, l} \in \{0, 1\} \quad \text{for } 1 \leq i \leq t.$$

The assumption $r_{\mathbb{R}}(\beta_{d_i, l})_{t \times k} = t$ implies that the number of solutions over $[1, n]$ is at most n^{k-t} . Finally, we have at most 2^{kt} possibilities on the left-hand side for $\alpha_{d_i, l}$ s and, similarly, we have at most 2^{kt} possibilities on the right-hand side for $\beta_{d_i, l}$ s, therefore the number of possible systems of equations is at most 2^{2k^2}

(2): The number of vectors $\mathbf{A}^{(k, n)}$ is $\binom{n}{k+1}$ as in Part 1. Fix a vector $\mathbf{A}^{(k, n)}$ and count the suitable vectors $\mathbf{B}^{(k, n)}$. It follows from the assumptions $r(\mathbf{H}_{\mathbf{A}^{(k, n)}}, \mathbf{H}_{\mathbf{B}^{(k, n)}}) = (r, r)$, $|\mathbf{H}_{\mathbf{A}^{(k)}} \cap \mathbf{H}_{\mathbf{B}^{(k)}}| = 2^r$ that the vectors $(\alpha_{d, 1}, \dots, \alpha_{d, k})$, $d = 1, \dots, 2^r$ and the vectors $(\beta_{d, 1}, \dots, \beta_{d, k})$, $d = 1, \dots, 2^r$, respectively form r -dimensional subspaces of \mathbb{F}_2^k . Considering the zero vectors of these subspaces we get $a_0 = b_0$. The integers b_1, \dots, b_k are solutions of the system of equations

$$a_0 + \sum_{l=1}^k \alpha_{d, l} a_l = b_0 + \sum_{l=1}^k \beta_{d, l} x_l \quad \alpha_{d, l}, \beta_{d, l} \in \{0, 1\} \quad \text{for } 1 \leq d \leq 2^r.$$

Similarly to the previous part this system of equation has at most n^{k-r} solutions over $[1, n]$ and the number of choices for the r linearly independent rows is at most 2^{2k^2} .

(3): Fix a vector $\mathbf{A}^{(k, n)}$. Let us suppose that for a vector $\mathbf{B}^{(k, n)}$ we have $r(\mathbf{H}_{\mathbf{A}^{(k, n)}}, \mathbf{H}_{\mathbf{B}^{(k, n)}}) = (k, k)$ and $|\mathbf{H}_{\mathbf{A}^{(k, n)}} \cap \mathbf{H}_{\mathbf{B}^{(k, n)}}| > 2^{k-1}$. Let the common vertices be

$$a_0 + \sum_{l=1}^k \alpha_{d, l} a_l = b_0 + \sum_{l=1}^k \beta_{d, l} b_l, \quad \alpha_{d, l}, \beta_{d, l} \in \{0, 1\} \quad \text{for } 1 \leq d \leq m,$$

where we may assume that the rows d_1, \dots, d_k are linearly independent, i.e. the matrix $B_k = (\beta_{d_i, l})_{k \times k}$ is regular. Write the rows d_1, \dots, d_k in matrix form as

$$(1) \quad \underline{a} = b_0 \underline{1} + B_k \underline{b},$$

with vectors $\underline{a} = (a_0 + \sum_{l=1}^k \alpha_{d_i, l} a_l)_{k \times 1}$, $\underline{1} = (1)_{k \times 1}$ and $\underline{b} = (b_i)_{k \times 1}$. It follows from (1) that

$$\underline{b} = B_k^{-1}(\underline{a} - b_0 \underline{1}) = B_k^{-1} \underline{a} - b_0 B_k^{-1} \underline{1}.$$

Let $B_k^{-1} \underline{1} = (d_i)_{k \times 1}$ and $B_k^{-1} \underline{a} = (c_i)_{k \times 1}$. Obviously, the number of subsets $\{i_1, \dots, i_t\} \subset \{1, \dots, k\}$ for which $d_{i_1} + \dots + d_{i_t} \neq 1$ is at least 2^{k-1} , therefore there exist $1 \leq u_1 < \dots < u_s \leq k$ and $1 \leq v_1 < \dots < v_t \leq k$ such that $a_0 + a_{u_1} + \dots + a_{u_s} = b_0 + b_{v_1} + \dots + b_{v_t}$, and $d_{v_1} + \dots + d_{v_t} \neq 1$. Hence

$$\begin{aligned} a_0 + a_{u_1} + \dots + a_{u_s} &= b_0 + b_{v_1} + \dots + b_{v_t} = b_0 + c_{v_1} + \dots + c_{v_t} - b_0(d_{v_1} + \dots + d_{v_t}) \\ b_0 &= \frac{a_0 + a_{u_1} + \dots + a_{u_s} - c_{v_1} - \dots - c_{v_t}}{1 - (d_{v_1} + \dots + d_{v_t})}. \end{aligned}$$

To conclude the proof we note that the number of sets $\{u_1, \dots, u_s\}$ and $\{v_1, \dots, v_t\}$ is at most 2^{2k} and there are at most 2^{k^2} choices for B_k and \underline{a} , respectively. Finally, for given B_k , \underline{a} , b_0 , $1 \leq u_1 < \dots < u_s \leq k$ and $1 \leq v_1 < \dots < v_t \leq k$, the vector $\mathbf{B}^{(k, n)}$ is determined uniquely. \square

In order to prove the theorems we need two lemmas from probability theory (see e.g. [1] p. 41, 95-98.). Let X_i be the indicator function of the event A_i and $S_n = X_1 + \dots + X_N$. For indices

i, j write $i \sim j$ if $i \neq j$ and the events A_i, A_j are dependant. We set $\Gamma = \sum_{i \sim j} \Pr\{A_i \cap A_j\}$ (the sum over ordered pairs).

Lemma 2.4. *If $E(S_n) \rightarrow \infty$ and $\Gamma = o(E(S_n)^2)$, then $S_n > 0$ a.e.*

In many instances, we would like to bound the probability that none of the bad events B_i , $i \in I$, occur. If the events are mutually independent, then $\Pr\{\cap_{i \in I} \overline{B_i}\} = \prod_{i \in I} \Pr\{\overline{B_i}\}$. When the B_i are "mostly" independent, the Janson's inequality allows us, sometimes, to say that these two quantities are "nearly" equal. Let Ω be a finite set and R be a random subset of Ω given by $\Pr\{r \in R\} = p_r$, these events being mutually independent over $r \in \Omega$. Let E_i , $i \in I$ be subsets of Ω , where I a finite index set. Let B_i be the event $E_i \subset R$. Let X_i be the indicator random variable for B_i and $X = \sum_{i \in I} X_i$ be the number of E_i s contained in R . The event $\cap_{i \in I} \overline{B_i}$ and $X = 0$ are then identical. For $i, j \in I$, we write $i \sim j$ if $i \neq j$ and $E_i \cap E_j \neq \emptyset$. We define $\Delta = \sum_{i \sim j} \Pr\{B_i \cap B_j\}$, here the sum is over ordered pairs. We set $M = \prod_{i \in I} \Pr\{\overline{B_i}\}$.

Lemma 2.5 (Janson's inequality). *Let $\varepsilon \in]0, 1[$, let $B_i, i \in I, \Delta, M$ be as above and assume that $\Pr\{B_i\} \leq \varepsilon$ for all i . Then*

$$M \leq \Pr\{\cap_{i \in I} \overline{B_i}\} \leq M e^{\frac{1}{1-\varepsilon} \frac{\Delta}{2}}.$$

Proof of Theorem 1.5. Let $\mathbf{H}_{\mathbf{A}_1^{(k,n)}}, \dots, \mathbf{H}_{\mathbf{A}_N^{(k,n)}}$ be the distinct non-degenerate k -cubes in $[1, n]$. Let B_i be the event $\mathbf{H}_{\mathbf{A}_i^{(k,n)}} \subset S(n, cn^{-\frac{k+1}{2k}})$. Then $\Pr\{B_i\} = c^{2k} n^{-(k+1)} = o(1)$ and $N = (1 + o(1)) \binom{n}{k+1} \frac{1}{k!}$. It is enough to prove

$$\Delta = \sum_{i \sim j} \Pr\{B_i \cap B_j\} = o(1)$$

since then Janson's inequality implies

$$\begin{aligned} \Pr\{S(n, cn^{-\frac{k+1}{2k}}) \text{ does not contain any } k\text{-cubes}\} &= \Pr\{\cap_{i=1}^N \overline{B_i}\} = \\ &= (1 + o(1)) (1 - (cn^{-\frac{k+1}{2k}})^{2k})^{(1+o(1)) \binom{n}{k+1} \frac{1}{k!}} = (1 + o(1)) e^{-\frac{c^{2k}}{(k+1)!k!}}. \end{aligned}$$

It remains to verify that $\sum_{i \sim j} \Pr\{B_i \cap B_j\} = o(1)$. We split this sum according to the ranks in the following way

$$\begin{aligned} \sum_{i \sim j} \Pr\{B_i \cap B_j\} &= \sum_{s=0}^k \sum_{t=0}^k \sum_{\substack{i \sim j \\ r(\mathbf{H}_{\mathbf{A}_i^{(k,n)}}, \mathbf{H}_{\mathbf{A}_j^{(k,n)}}) = (s,t)}} \Pr\{B_i \cap B_j\} = \\ &= 2 \sum_{s=1}^k \sum_{t=0}^{s-1} \sum_{\substack{i \sim j \\ r(\mathbf{H}_{\mathbf{A}_i^{(k,n)}}, \mathbf{H}_{\mathbf{A}_j^{(k,n)}}) = (s,t)}} \Pr\{B_i \cap B_j\} + \\ &= \sum_{r=0}^{k-1} \sum_{\substack{i \sim j \\ r(\mathbf{H}_{\mathbf{A}_i^{(k,n)}}, \mathbf{H}_{\mathbf{A}_j^{(k,n)}}) = (r,r) \\ |\mathbf{H}_{\mathbf{A}_i^{(k,n)}} \cap \mathbf{H}_{\mathbf{A}_j^{(k,n)}}| = 2^r}} \Pr\{B_i \cap B_j\} + \sum_{r=1}^{k-1} \sum_{\substack{i \sim j \\ r(\mathbf{H}_{\mathbf{A}_i^{(k,n)}}, \mathbf{H}_{\mathbf{A}_j^{(k,n)}}) = (r,r) \\ |\mathbf{H}_{\mathbf{A}_i^{(k,n)}} \cap \mathbf{H}_{\mathbf{A}_j^{(k,n)}}| < 2^r}} \Pr\{B_i \cap B_j\} + \end{aligned}$$

$$\sum_{\substack{i \sim j \\ r(\mathbf{H}_{A_i^{(k,n)}}, \mathbf{H}_{A_j^{(k,n)}}) = (k, k) \\ |\mathbf{H}_{A_i^{(k,n)}} \cap \mathbf{H}_{A_j^{(k,n)}}| \leq 2^{k-1}}} \Pr\{B_i \cap B_j\} + \sum_{\substack{i \sim j \\ r(\mathbf{H}_{A_i^{(k,n)}}, \mathbf{H}_{A_j^{(k,n)}}) = (k, k) \\ |\mathbf{H}_{A_i^{(k,n)}} \cap \mathbf{H}_{A_j^{(k,n)}}| > 2^{k-1}}} \Pr\{B_i \cap B_j\}.$$

The first sum can be estimated by Lemmas 2.2 and 2.3 (3)

$$\begin{aligned} \sum_{s=1}^k \sum_{t=0}^{s-1} \sum_{\substack{i \sim j \\ r(\mathbf{H}_{A_i^{(k,n)}}, \mathbf{H}_{A_j^{(k,n)}}) = (s, t) \\ |\mathbf{H}_{A_i^{(k,n)}} \cap \mathbf{H}_{A_j^{(k,n)}}| \leq 2^{k-1}}} \Pr\{B_i \cap B_j\} &\leq \sum_{s=1}^k \sum_{t=0}^{s-1} 2^{2k^2} \binom{n}{k+1} n^{k+1-s} \left(\frac{c}{n^{\frac{k+1}{2^k}}} \right)^{2 \cdot 2^{k-2t}} = \\ &n^{o(1)} \sum_{s=1}^k \frac{n^{2^{s-1} \frac{k+1}{2^k}}}{n^s} = n^{o(1)} (n^{\frac{k+1}{2^k} - 1} + n^{\frac{k+1}{2} - k}) = o(1), \end{aligned}$$

since the sequence $a_s = 2^{s-1} \frac{k+1}{2^k} - s$ is decreasing for $1 \leq s \leq k - \log_2(k+1) + 1$ and increasing for $k - \log_2(k+1) + 1 < s \leq k$.

To estimate the second sum we apply Lemma 2.3 (2)

$$\begin{aligned} \sum_{r=0}^{k-1} \sum_{\substack{i \sim j \\ r(\mathbf{H}_{A_i^{(k,n)}}, \mathbf{H}_{A_j^{(k,n)}}) = (r, r) \\ |\mathbf{H}_{A_i^{(k,n)}} \cap \mathbf{H}_{A_j^{(k,n)}}| = 2^r}} \Pr\{B_i \cap B_j\} &\leq \sum_{r=0}^{k-1} 2^{2k^2} \binom{n}{k+1} n^{k-r} \left(\frac{c}{n^{\frac{k+1}{2^k}}} \right)^{2 \cdot 2^{k-2r}} = \\ &n^{-1+o(1)} \sum_{r=0}^{k-1} \frac{n^{2^r \frac{k+1}{2^k}}}{n^r} = n^{-1+o(1)} (n^{\frac{k+1}{2^k}} + n^{\frac{k+1}{2} - (k-1)}) = o(1). \end{aligned}$$

The third sum can be bounded using Lemma 2.3 (1)

$$\begin{aligned} \sum_{r=1}^{k-1} \sum_{\substack{i \sim j \\ r(\mathbf{H}_{A_i^{(k,n)}}, \mathbf{H}_{A_j^{(k,n)}}) = (r, r) \\ |\mathbf{H}_{A_i^{(k,n)}} \cap \mathbf{H}_{A_j^{(k,n)}}| < 2^r}} \Pr\{B_i \cap B_j\} &\leq \sum_{r=1}^{k-1} 2^{2k^2} \binom{n}{k+1} n^{k+1-r} \left(\frac{c}{n^{\frac{k+1}{2^k}}} \right)^{2 \cdot 2^{k-2r+1}} \leq \\ &n^{o(1) - \frac{k+1}{2^k}} \sum_{r=1}^{k-1} \frac{n^{2^r \frac{k+1}{2^k}}}{n^r} = n^{o(1) - \frac{k+1}{2^k}} (n^{2^{\frac{k+1}{2^k}} - 1} + n^{\frac{k+1}{2} - (k-1)}) = o(1). \end{aligned}$$

Similarly, for the fourth sum we apply Lemma 2.3 (1)

$$\sum_{\substack{i \sim j \\ r(\mathbf{H}_{A_i^{(k,n)}}, \mathbf{H}_{A_j^{(k,n)}}) = (k, k) \\ |\mathbf{H}_{A_i^{(k,n)}} \cap \mathbf{H}_{A_j^{(k,n)}}| \leq 2^{k-1}}} \Pr\{B_i \cap B_j\} \leq n^{o(1)} n^{k+2} \left(\frac{c}{n^{\frac{k+1}{2^k}}} \right)^{1.5 \cdot 2^k} = o(1).$$

To estimate the fifth sum we note that $|\mathbf{H}_{\mathbf{A}_i^{(k,n)}} \cup \mathbf{H}_{\mathbf{A}_j^{(k,n)}}| \geq 2^k + 1$. It follows from Lemma 3.3 that

$$\sum_{\substack{i \sim j \\ r(\mathbf{H}_{\mathbf{A}_i^{(k,n)}}, \mathbf{H}_{\mathbf{A}_j^{(k,n)}}) = (k, k) \\ |\mathbf{H}_{\mathbf{A}_i^{(k,n)}} \cap \mathbf{H}_{\mathbf{A}_j^{(k,n)}}| > 2^{k-1}}} \Pr\{B_i \cap B_j\} \leq 2^{2k^2+2k} n^{k+1} \left(\frac{c}{n^{2^k}}\right)^{2^k+1} = o(1),$$

which completes the proof. \square

Proof of Theorem 1.6. Let $\epsilon > 0$ and for simplicity let $D_n = D_n(\epsilon)$ and $E_n = E_n(\epsilon)$. In the proof we use the estimations

$$(2) \quad 2^{2D_n} \leq 2^{2^{\log_2 \log_2 n + \log_2 \log_2 \log_2 n + \frac{(1-\epsilon) \log_2 \log_2 \log_2 n}{\log 2 \log_2 \log_2 n}}} = n^{\log_2 \log_2 n + (1-\epsilon+o(1)) \log_2 \log_2 \log_2 n}$$

and

$$(3) \quad 2^{2E_n+1} \geq 2^{2^{\log_2 \log_2 n + \log_2 \log_2 \log_2 n + \frac{(1+\epsilon) \log_2 \log_2 \log_2 n}{\log 2 \log_2 \log_2 n}}} = n^{\log_2 \log_2 n + (1+\epsilon+o(1)) \log_2 \log_2 \log_2 n}$$

In order to verify Theorem 2 we have to show that

$$(4) \quad \lim_{n \rightarrow \infty} \Pr\{S(n, \frac{1}{2}) \text{ contains a } D_n\text{-cube}\} = 1$$

and

$$(5) \quad \lim_{n \rightarrow \infty} \Pr\{S(n, \frac{1}{2}) \text{ contains an } (E_n + 1)\text{-cube}\} = 0.$$

To prove the limit in (4) let $\mathbf{H}_{\mathbf{A}_1^{(D_n,n)}}, \dots, \mathbf{H}_{\mathbf{A}_N^{(D_n,n)}}$ be the different non-degenerate D_n -cubes in $[1, n]$, B_i be the event $H_{\mathbf{A}_i^{(D_n,n)}} \subset S(n, \frac{1}{2})$, X_i be the indicator random variable for B_i and $X = X_1 + \dots + X_N$ be the number of $\mathbf{H}_{\mathbf{A}_i^{(D_n,n)}} \subset S(n, \frac{1}{2})$. The linearity of expectation gives by Lemma 1 and inequality (2)

$$E(X) = NE(X_i) = (1 + o(1)) \binom{n}{D_n + 1} \frac{1}{D_n!} 2^{-2D_n} \geq$$

$$n^{\log_2 \log_2 n + (1+o(1)) \log_2 \log_2 \log_2 n} n^{-\log_2 \log_2 n - (1-\epsilon+o(1)) \log_2 \log_2 \log_2 n} = n^{(\epsilon+o(1)) \log_2 \log_2 \log_2 n},$$

therefore $E(X) \rightarrow \infty$, as $n \rightarrow \infty$. By Lemma 2.4 it remains to prove that

$$\sum_{i \sim j} \Pr\{B_i \cap B_j\} = o(E(X)^2)$$

where $i \sim j$ means that the events B_i, B_j are not independent i.e. the cubes $\mathbf{H}_{\mathbf{A}_i^{(D_n,n)}}, \mathbf{H}_{\mathbf{A}_j^{(D_n,n)}}$ have common vertices. We split this sum according to the ranks

$$(6) \quad \sum_{i \sim j} \Pr\{B_i \cap B_j\} = \sum_{s=0}^{D_n} \sum_{t=0}^{D_n} \sum_{\substack{i \sim j \\ r(\mathbf{H}_{\mathbf{A}_i^{(D_n,n)}}, \mathbf{H}_{\mathbf{A}_j^{(D_n,n)}}) = (s, t)}} \Pr\{B_i \cap B_j\} \leq$$

$$\sum_{\substack{i \sim j \\ r(\mathbf{H}_{\mathbf{A}_i^{(D_n,n)}}, \mathbf{H}_{\mathbf{A}_j^{(D_n,n)}}) = (0, 0)}} \Pr\{B_i \cap B_j\} + 2 \sum_{s=1}^{D_n} \sum_{t=0}^s \sum_{\substack{i \sim j \\ r(\mathbf{H}_{\mathbf{A}_i^{(D_n,n)}}, \mathbf{H}_{\mathbf{A}_j^{(D_n,n)}}) = (s, t)}} \Pr\{B_i \cap B_j\}.$$

The condition $r(\mathbf{H}_{\mathbf{A}_i^{(D_n, n)}}, \mathbf{H}_{\mathbf{A}_j^{(D_n, n)}}) = (0, 0)$ implies that $|\mathbf{H}_{\mathbf{A}_i^{(D_n, n)}} \cup \mathbf{H}_{\mathbf{A}_j^{(D_n, n)}}| = 2^{D_n+1} - 1$, thus by Lemma 3.2

$$\sum_{\substack{i \sim j \\ r(\mathbf{H}_{\mathbf{A}_i^{(D_n, n)}}, \mathbf{H}_{\mathbf{A}_j^{(D_n, n)}}) = (0, 0)}} \Pr\{B_i \cap B_j\} \leq 2^{2D_n^2} \binom{n}{D_n+1} n^{D_n} 2^{-2^{D_n+1}+1} = o\left(\left(\binom{n}{D_n+1} \frac{1}{D_n!} 2^{-2^{D_n}}\right)^2\right) = o(E(X)^2).$$

In the light of Lemmas 2.2 and 2.3 (1) the second term in (6) can be estimated as

$$\sum_{s=1}^{D_n} \sum_{t=0}^s \sum_{i \sim j} \Pr\{B_i \cap B_j\} \leq \sum_{s=1}^{D_n} \sum_{t=0}^s \binom{n}{D_n+1} 2^{2D_n^2} n^{D_n+1-s} 2^{-2 \cdot 2^{D_n+2t}} = \left(\binom{n}{D_n+1} \frac{1}{D_n!} 2^{-2^{D_n}}\right)^2 n^{o(1)} \sum_{s=1}^{D_n} \sum_{t=0}^s \frac{2^{2t}}{n^s} = \left(\binom{n}{D_n+1} \frac{1}{D_n!} 2^{-2^{D_n}}\right)^2 n^{o(1)} \sum_{s=1}^{D_n} \frac{2^{2s}}{n^s}$$

Finally, the function $f(x) = \frac{2^{2x}}{n^x}$ decreases on $(-\infty, \log_2 \log n - 2 \log_2 \log 2]$ and increases on $[\log_2 \log n - 2 \log_2 \log 2, \infty)$, therefore by (2)

$$\sum_{s=1}^{D_n} \frac{2^{2s}}{n^s} = n^{o(1)} \left(\frac{4}{n} + \frac{2^{2^{D_n}}}{n^{D_n}}\right) = n^{-1+o(1)},$$

which proves the limit in (4).

In order to prove the limit in (5) let $\mathbf{H}_{\mathbf{C}_1^{(E_n+1, n)}}, \dots, \mathbf{H}_{\mathbf{C}_K^{(E_n+1, n)}}$ be the distinct (E_n+1) -cubes in $[1, n]$ and let F_i be the event $\mathbf{H}_{\mathbf{C}_i^{(E_n+1, n)}} \subset S(n, \frac{1}{2})$. By (3) we have

$$\Pr\{S_n \text{ contains an } (E_n+1)\text{-cube}\} = \Pr\{\cup_{i=1}^K F_i\} \leq \sum_{i=1}^K \Pr\{F_i\} \leq \binom{n}{E_n+2} 2^{-2^{E_n+1}} \leq \frac{n^{\log_2 \log_2 n + (1+o(1)) \log_2 \log_2 \log_2 n}}{n^{\log_2 \log_2 n + (1+\epsilon+o(1)) \log_2 \log_2 \log_2 n}} = o(1),$$

which completes the proof. \square

3. Concluding remarks

The aim of this paper is to study non-degenerate Hilbert cubes in a random sequence. A natural problem would be to give analogous theorems for Hilbert cubes, where degenerate cubes are allowed. In this situation the dominant terms may come from arithmetic progressions. An AP_{k+1} forms a k -cube. One can prove by the Janson inequality (see Lemma 2.5) that for a fixed $k \geq 2$

$$\lim_{n \rightarrow \infty} \Pr\{S(n, cn^{-\frac{2}{k+1}}) \text{ contains no } AP_{k+1}\} = e^{-\frac{c^{k+1}}{2^k}}.$$

An easy argument shows (using Janson's inequality again) that for all $c > 0$, with $p_n = cn^{-2/5}$

$$\lim_{n \rightarrow \infty} \Pr\{S(n, p_n) \text{ contains no 4-cubes}\} = e^{-\frac{c^5}{8}}.$$

Conjecture 3.1. For $k \geq 4$

$$\lim_{n \rightarrow \infty} \Pr\{S(n, cn^{-\frac{2}{k+1}}) \text{ contains no } k\text{-cubes}\} = e^{-\frac{c^{k+1}}{2k}}.$$

A simple calculation implies that in the random subset $S(n, 1/2)$ the length of the longest arithmetic progression is a.e. nearly $2 \log_2 n$, therefore it contains a Hilbert cube of dimension $(2 - \varepsilon) \log_2 n$.

Conjecture 3.2. For every $\varepsilon > 0$

$$\lim_{n \rightarrow \infty} \Pr\{\text{the maximal dimension of Hilbert cubes in } S(n, \frac{1}{2}) \text{ is } < (2 + \varepsilon) \log_2 n\} = 1.$$

N. Hegyvári (see [5]) studied the special case where the generating elements of Hilbert cubes are distinct. He proved that in this situation the maximal dimension of Hilbert cubes is a.e. between $c_1 \log n$ and $c_2 \log n \log \log n$. In this problem the lower bound seems to be the correct magnitude.

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Csaba SÁNDOR
 Department of Stochastics
 Budapest University of Technology and Economics
 Egrý J. u. 1, 1111 Budapest
 Hungary

E-mail : csandor@math.bme.hu