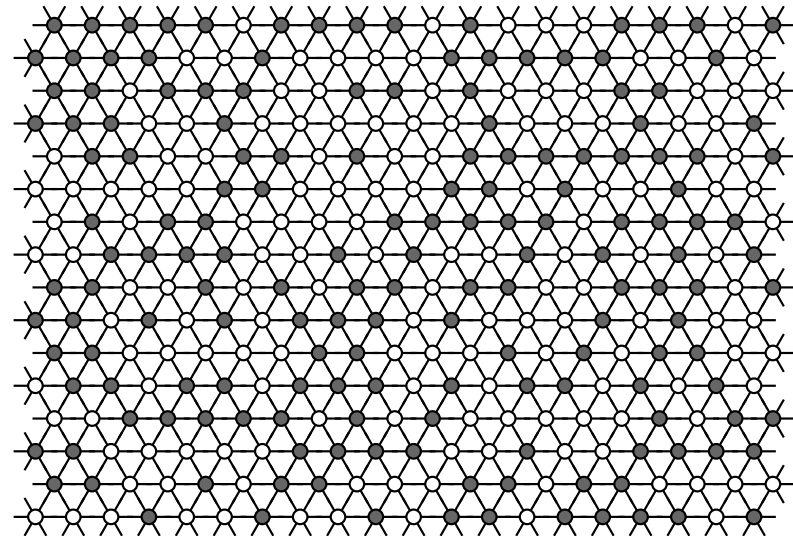
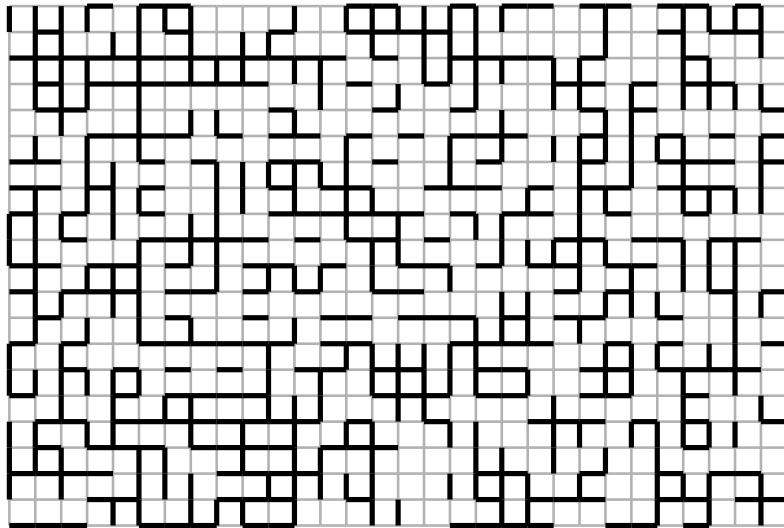


Noise and dynamical sensitivity in critical planar percolation

Based on C. Garban, G. Pete & O. Schramm:
The Fourier spectrum of critical percolation, *Acta Math.* 2010.

Bernoulli(p) bond and site percolation

Given an (infinite) graph $G = (V, E)$ and $p \in [0, 1]$. Each site (or bond) is chosen open with probability p , closed with $1 - p$, independently of each other. Consider the **open connected clusters**. $\theta(p) := \mathbf{P}_p[0 \longleftrightarrow \infty]$.



Theorem (Harris 1960 and Kesten 1980).

$$p_c(\mathbb{Z}^2, \text{bond}) = p_c(\Delta, \text{site}) = 1/2, \text{ and } \theta(1/2) = 0.$$

For $p > 1/2$, there is a.s. one infinite cluster.

Crossing probabilities and criticality

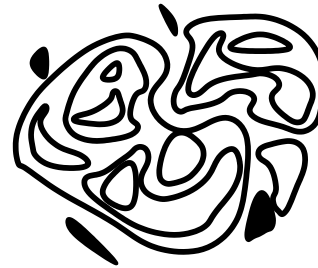
$p \approx 0.8$



$p \approx 0.55$



$p = 0.5$



$p \approx 0.45$



Theorem (Russo 1978 and Seymour-Welsh 1978). In critical percolation on almost any planar lattice, for $L, n > 0$,

$$0 < a_L < \mathbf{P}[\text{left-right crossing in } n \times Ln] < b_L < 1.$$

Same holds for annulus-crossings.

By repeating this on all scales, and gluing the pieces by FKG:

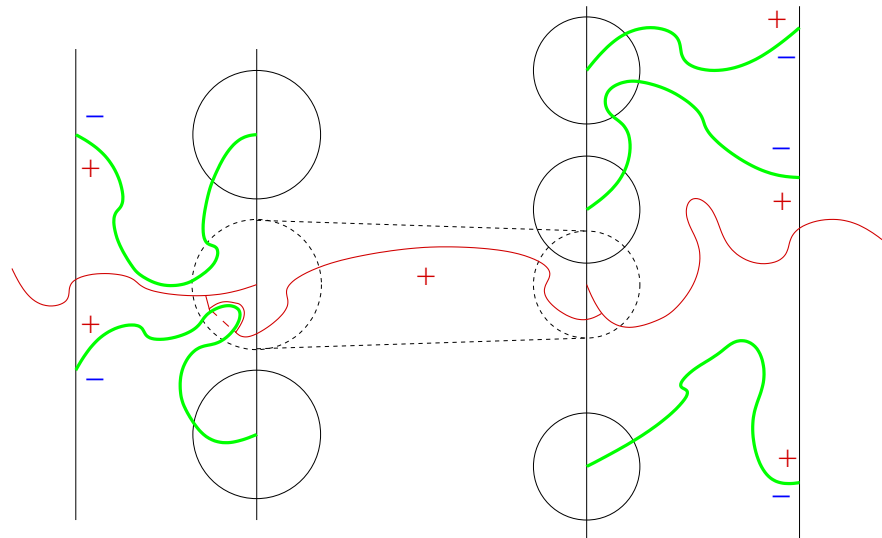
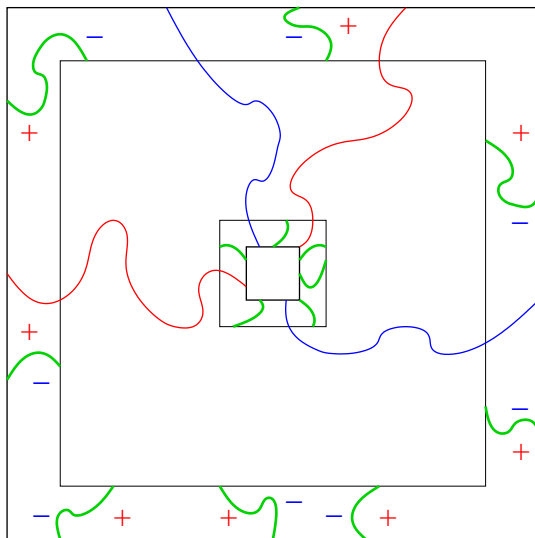
$$(r/R)^\alpha < \mathbf{P}[\partial B_r \longleftrightarrow \partial B_R] < (r/R)^\beta.$$

Moreover, for the (polychromatic) ℓ -arm probabilities

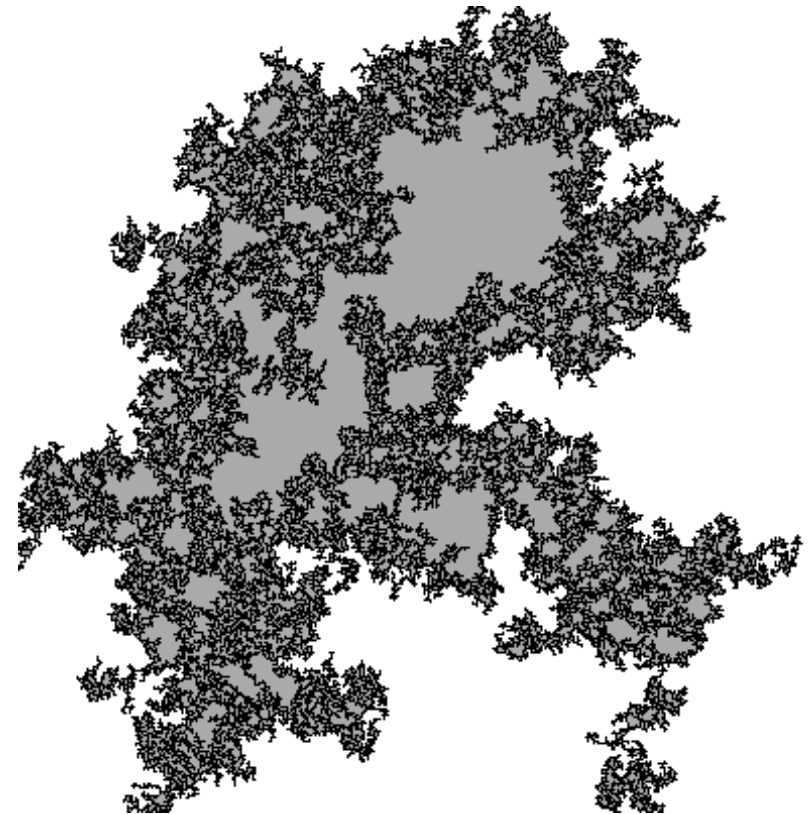
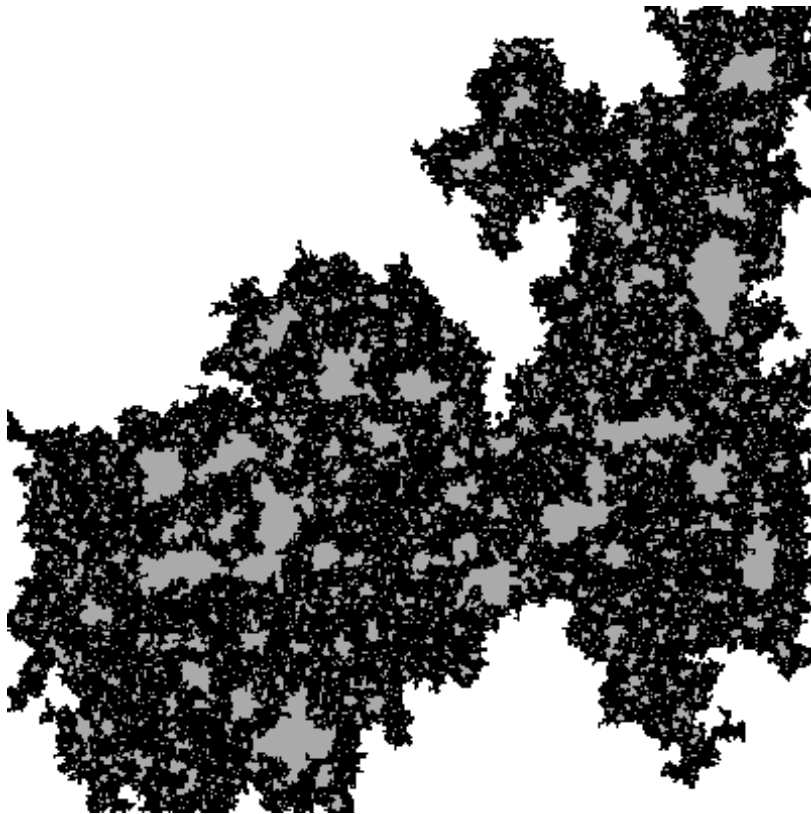
$$\alpha_\ell(r, R) := \mathbf{P}[\partial B_r \xleftrightarrow{\ell} \partial B_R],$$

again have **quasi-multiplicativity**: $\alpha_\ell(r, R) \asymp \alpha_\ell(r, \rho) \alpha_\ell(\rho, R)$, and thus $c_\ell (r/R)^{C_\ell} < \alpha_\ell(r, R) < C_\ell (r/R)^{c_\ell}$.

But these are non-monotone events, so cannot use just FKG to prove this q-multiplicativity. Need **Separation Lemma**: conditioned on having ℓ arms from ∂B_r to ∂B_R , the collection of interfaces both in $B_R \setminus B_{R/2}$ and $B_{2r} \setminus B_r$ are well-separated. Then we can glue.



Bernoulli(1/2) bond and site percolation

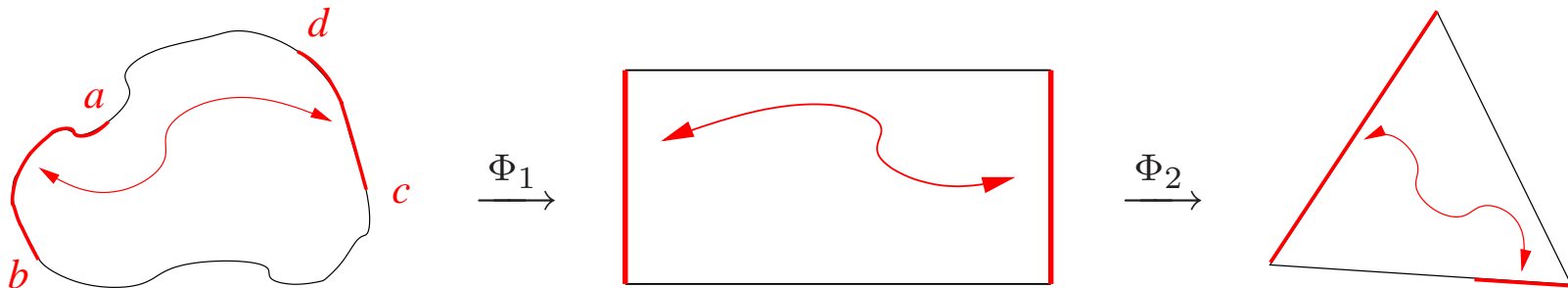


Conformal invariance on Δ

Theorem (Smirnov 2001). For $p = 1/2$ site percolation on Δ_η , and $Q \subset \mathbb{C}$ a piecewise smooth quad (simply connected domain with four boundary points $\{a, b, c, d\}$),

$$\lim_{\eta \rightarrow 0} \mathbf{P} \left[ab \longleftrightarrow cd \text{ inside } Q, \text{ in percolation on } \Delta_\eta \right]$$

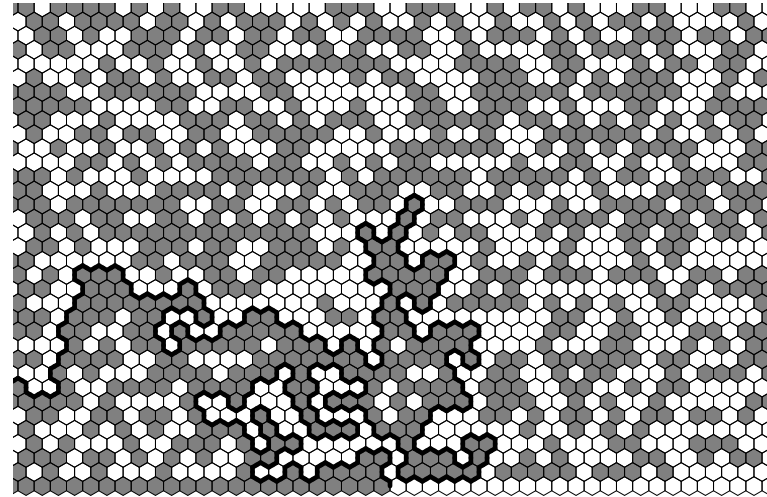
exists, is strictly between 0 and 1, and conformally invariant.



Calls for a **continuum scaling limit**, encoding macroscopic connectivity, cluster boundaries, etc. **Aizenman '95**, **Schramm '00**, **Camia-Newman '06**, **Sheffield '09**, **Schramm-Smirnov '10**. In physics, correlation functions.

SLE_6 exponents

Given the conformal invariance, the exploration path converges to the **Stochastic Loewner Evolution** with $\kappa = 6$ (Schramm 2000).



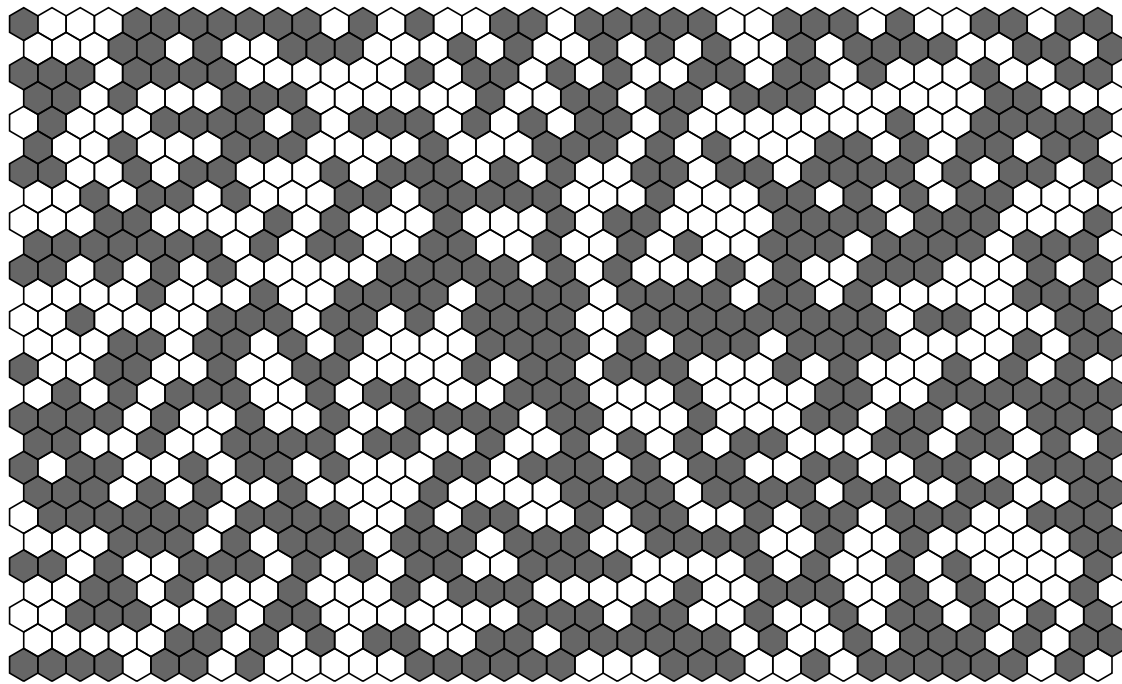
Using the SLE_6 curve, several **critical exponents** can be computed (Lawler-Schramm-Werner, Smirnov-Werner 2001, plus Kesten 1987), e.g.:

$$\alpha_4(r, R) := \mathbf{P} \left[\begin{array}{c} R \\ \text{Diagram of a circle of radius R with a smaller circle of radius r inside. Two paths, one red and one blue, start from the center and end on the boundary.} \\ r \end{array} \right] = (r/R)^{5/4+o(1)},$$

$\alpha_1(r, R) = (r/R)^{5/48+o(1)}$, and $\theta(p_c + \epsilon) := \mathbf{P}_{p_c+\epsilon}[0 \longleftrightarrow \infty] = \epsilon^{5/36+o(1)}$.
 Here $\beta = 5/36 = \frac{5/48}{2-5/4} = \frac{\xi_1}{2-\xi_4}$. (Will explain this numerology tomorrow.)

Percolation and noise

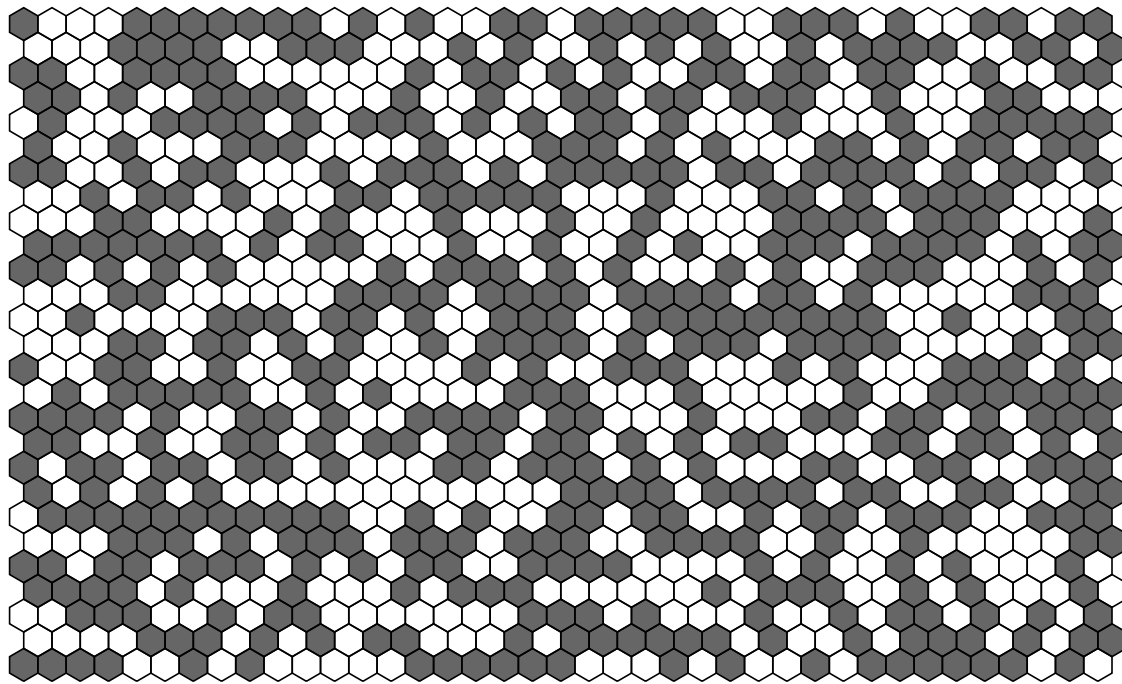
Take an ω critical percolation configuration. Let ω^ϵ be a new configuration, where each site (or bond) is **resampled** with probability ϵ , independently. (The ϵ -noised version of ω .)



For how large an ϵ can we still predict from ω whether there is a left-right crossing in ω^ϵ ?

Percolation and noise

Take an ω critical percolation configuration. Let ω^ϵ be a new configuration, where each site (or bond) is **resampled** with probability ϵ , independently. (The ϵ -noised version of ω .)



For how large an ϵ can we still predict from ω whether there is a left-right crossing in ω^ϵ ?

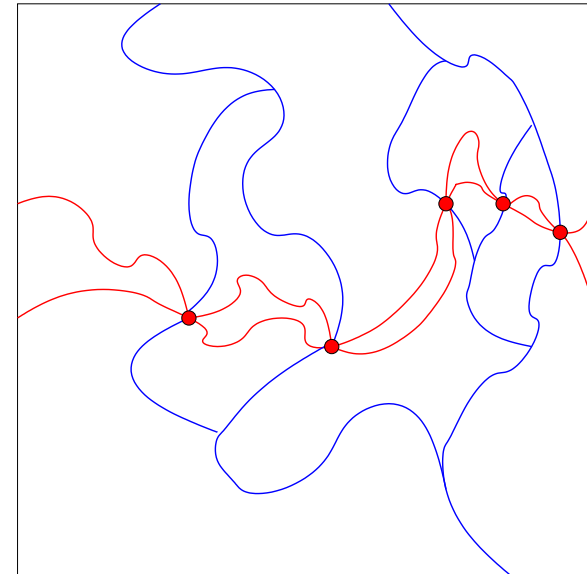
Naive idea: how many pivotals are there?

A site (or bond) is **pivotal** in ω , if flipping it changes the existence of a left-right crossing.

$$\mathbf{E}|\mathbf{Piv}_n| \asymp n^2 \alpha_4(n) \quad (= n^{3/4+o(1)}).$$

Furthermore, $\mathbf{E}[|\mathbf{Piv}_n|^2] \leq C (\mathbf{E}|\mathbf{Piv}_n|)^2$.
So, $\mathbf{P}[|\mathbf{Piv}_n| > \lambda \mathbf{E}|\mathbf{Piv}_n|] < C/\lambda^2$, any λ .

And not only $\exists \epsilon \mathbf{P}[|\mathbf{Piv}_n| > \epsilon \mathbf{E}|\mathbf{Piv}_n|] > \epsilon$,
but $\mathbf{P}[0 < |\mathbf{Piv}_n| < \epsilon \mathbf{E}|\mathbf{Piv}_n|] \asymp \epsilon^{11/9+o(1)}$,
as $\epsilon \rightarrow 0$ (exponent only for Δ).



Cannot have many pivotals \implies If $\epsilon_n \mathbf{E}[|\mathbf{Piv}_n|] \rightarrow 0$, then we don't hit any pivotals (even in expectation) \implies Asymptotically full correlation.

Cannot have few pivotals (if there's any) \implies If $\epsilon_n \mathbf{E}[|\mathbf{Piv}_n|] \rightarrow \infty$, hit many pivotals (at least in expectation). But $\not\implies$ asymptotic independence!

Noise sensitivity of percolation

All results use Fourier analysis of Boolean functions:

Theorem (Benjamini, Kalai & Schramm 1998). If $\epsilon > 0$ is fixed, and f_n is the indicator function for a left-right percolation crossing in an $n \times n$ square, then as $n \rightarrow \infty$

$$\mathbf{E}[f_n(\omega) f_n(\omega^\epsilon)] - \mathbf{E}[f_n(\omega)]^2 \rightarrow 0.$$

This holds for all $\epsilon = \epsilon_n > c/\log n$.

Theorem (Schramm & Steif 2005). Same if $\epsilon_n > n^{-a}$ for some positive $a > 0$. If triangular lattice, may take any $a < 1/8$.

Theorem (Garban, P & Schramm 2008). Same holds if and only if $\epsilon_n \mathbf{E}[|\text{pivotal}|] \rightarrow \infty$. For triangular lattice, this threshold is $\epsilon_n = n^{-3/4+o(1)}$.

Dynamical percolation

Each variable is resampled according to an independent Poisson(1) clock. This is a Markov process $\{\omega(t) : t \in [0, \infty)\}$, in which $\omega(t + s)$ is an ϵ -noised version of $\omega(t)$, with $\epsilon = 1 - \exp(-s)$.

An **exceptional time** is such a (random) t , at which an almost sure property of the static process fails for $\omega(t)$.

Main example: (Non-)existence of an infinite cluster in percolation.

Toy example: Brownian motion on the circle does sometimes hit a given point, as opposed to its static version: a uniform random point.

In this toy example, the set of exceptional times is a random Cantor set of Lebesgue measure zero (because of Fubini) and Hausdorff-dimension $1/2$.

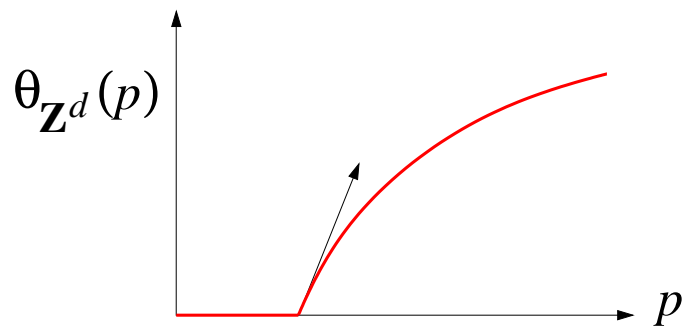
If the static event is **not extremely unlikely**, and it is **very sensitive to noise**, then we may have some chance to see an exceptional time.

Dynamical percolation results

Theorem (Häggström, Peres & Steif 1997).

- No exceptional times when $p \neq p_c$.
- No exceptional times when $p = p_c$ for bond percolation on \mathbb{Z}^d , $d \geq 19$.

The latter is essentially due to Hara-Slade '90 on the off-critical exponent $\beta = 1$:



$\theta(p_c + \epsilon) < C\epsilon$, hence, even switching asymmetrically, $\mathbf{E}[\text{number of } \epsilon\text{-subintervals of } [0, 1] \text{ with } 0 \longleftrightarrow \infty] \leq C$. But this exceptional set is closed without isolated points, so this number should blow up, if non-zero.

Theorem (Schramm & Steif 2005).

- There are exceptional times (a.s.) for critical site percolation on the triangular lattice.
- They have Hausdorff dimension in $[1/6, 31/36]$.

Theorem (Garban, P & Schramm 2008).

- There are exceptional times also on \mathbb{Z}^2 .
- On the triangular grid they have Hausdorff dimension $31/36$.
- On the triangular grid, there are exceptional times with an infinite white **and** an infinite black cluster simultaneously. ($1/9 \leq \dim \leq 2/3$)

What is the Fourier spectrum and why is it useful?

$f_n : \{\pm 1\}^{V_n} \longrightarrow \{\pm 1\}$ indicator of left-right crossing, $V = V_n$ vertices.

$(N_\epsilon f)(\omega) := \mathbf{E}[f(\omega^\epsilon) \mid \omega]$ is the **noise operator**, acting on the space $L^2(\Omega, \mu)$, where $\Omega = \{\pm 1\}^V$, μ uniform measure, inner product $\mathbf{E}[fg]$.

Correlation: $\mathbf{E}[f(\omega^\epsilon)f(\omega)] - \mathbf{E}[f(\omega)]\mathbf{E}[f(\omega^\epsilon)] = \mathbf{E}[f(\omega)N_\epsilon f(\omega)] - \mathbf{E}[f(\omega)]^2$. So, we would like to **diagonalize** the noise operator N_ϵ .

Let χ_i be the function $\chi_i(\omega) = \omega(i)$, $\omega \in \Omega$.

For $S \subset V$, let $\chi_S := \prod_{i \in S} \chi_i$, the **parity inside S** . Then

$$N_\epsilon \chi_i = (1 - \epsilon) \chi_i; \quad N_\epsilon \chi_S = (1 - \epsilon)^{|S|} \chi_S.$$

Moreover, the family $\{\chi_S, S \subseteq V\}$ is an **orthonormal basis** of $L^2(\Omega, \mu)$.

Any function $f \in L^2(\Omega, \mu)$ in this basis (**Fourier-Walsh series**):

$$\hat{f}(S) := \mathbf{E}[f\chi_S]; \quad f = \sum_{S \subseteq V} \hat{f}(S) \chi_S.$$

The correlation:

$$\begin{aligned} \mathbf{E}[fN_\epsilon f] - \mathbf{E}[f]^2 &= \sum_S \sum_{S'} \hat{f}(S) \hat{f}(S') \mathbf{E}[\chi_S N_\epsilon \chi_{S'}] - \mathbf{E}[f\chi_\emptyset]^2 \\ &= \sum_{\emptyset \neq S \subseteq V} \hat{f}(S)^2 (1 - \epsilon)^{|S|} = \sum_{k=1}^{|V_n|} (1 - \epsilon)^k \sum_{|S|=k} \hat{f}(S)^2. \end{aligned}$$

By Parseval, $\sum_S \hat{f}(S)^2 = \mathbf{E}[f^2] = 1$. So can define probability measure $\mathbf{P}[\mathcal{S}_f = S] := \hat{f}(S)^2 / \mathbf{E}[f^2]$, the **spectral sample** $\mathcal{S}_f \subseteq V$.

If, for some functions f_n and numbers k_n , we have $\mathbf{P}[0 < |\mathcal{S}_n| < tk_n] \rightarrow 0$ as $t \rightarrow 0$, uniformly in n , then $(1 - \epsilon)^k \sim \exp(-\epsilon k)$ implies that for $\epsilon_n \gg 1/k_n$ we have **asymptotic independence**. Maybe with $k_n = \mathbf{E}|\mathcal{S}_n|$?

Pivotals versus spectral sample

$\nabla_i f(\omega) := f(\sigma_i(\omega)) - f(\omega) \in \{-2, 0, +2\}$ gradient.

$\nabla_i f(\omega) = \sum_S \hat{f}(S) [\chi_S(\sigma_i(\omega)) - \chi_S(\omega)]$, hence $\widehat{\nabla_i f}(S) = -2\hat{f}(S)\mathbf{1}_{i \in S}$.

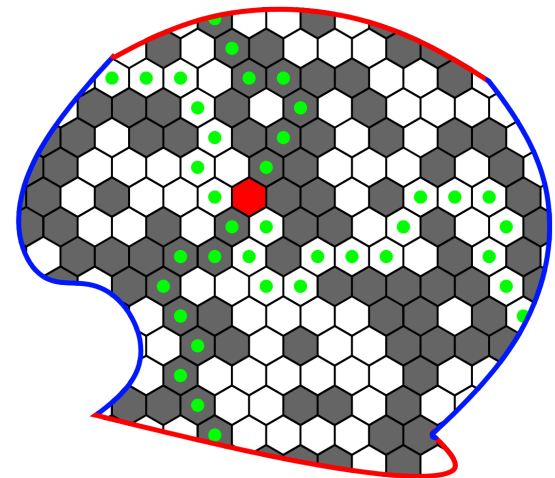
$\mathbf{P}[i \in \text{Piv}_f] = \frac{1}{4} \|\nabla_i f\|_2^2 = \frac{1}{4} \sum_S \widehat{\nabla_i f}(S)^2 = \sum_{S \ni i} \hat{f}(S)^2 = \mathbf{P}[i \in \mathcal{S}_f]$.

Thus, $\mathbf{E}|\mathcal{S}_f| = \mathbf{E}|\text{Piv}_f|$. So, the pivotal upper bound for noise sensitivity is sharp if there is tightness around $\mathbf{E}|\mathcal{S}|$.

Will see $\mathbf{P}[i, j \in \text{Piv}_f] = \mathbf{P}[i, j \in \mathcal{S}_f]$, hence $\mathbf{E}|\mathcal{S}_f|^2 = \mathbf{E}|\text{Piv}_f|^2$.

Not for more points and higher moments!
Both random subsets measure the “influence” or “relevance” of bits, but in different ways.

For percolation, $\mathbf{E}[|\text{Piv}_n|^2] \leq C (\mathbf{E}|\text{Piv}_n|)^2$,
hence $\exists c > 0$ s.t. $\mathbf{P}[|\mathcal{S}_n| > c\mathbf{E}|\mathcal{S}_n|] > c$.
That’s why one hopes for tightness around mean.



Three very simple examples

Dictator $_n(x_1, \dots, x_n) := x_1$.

Here $\text{Cov}[\text{Dic}_n(x), \text{Dic}_n(x^\epsilon)] = 1 - \epsilon$, so noise-stable.

And $\mathbf{P}[\mathcal{S}_n = \{x_1\}] = 1$.

Majority $_n(x_1, \dots, x_n) := \text{sgn}(x_1 + \dots + x_n) \approx \frac{1}{\sqrt{n}}(x_1 + \dots + x_n)$.

Here $\text{Cov}[\text{Maj}_n(x), \text{Maj}_n(x^\epsilon)] = 1 - O(\epsilon)$, so noise-stable.

And $\mathbf{P}[\mathcal{S}_n = \{x_i\}] \asymp 1/n$, most of the weight is on singletons.

On the other hand, $\mathbf{E}|\mathcal{S}_n| = \mathbf{E}|\text{Piv}_n| \asymp \frac{1}{\sqrt{n}} n \asymp \sqrt{n}$.

Parity $_n(x_1, \dots, x_n) := x_1 \cdots x_n$

Here $\text{Cov}[\text{Par}_n(x), \text{Par}_n(x^\epsilon)] = (1 - \epsilon)^n$, the most sensitive to noise.

And $\mathbf{P}[\mathcal{S}_n = \{x_1, \dots, x_n\}] = 1$.

Benjamini, Kalai & Schramm 1998

Theorem. A sequence f_n of monotone Boolean functions is **noise sensitive**, i.e., for any fixed $\epsilon > 0$,

$$\mathbf{E}[f_n(\omega) f_n(\omega^\epsilon)] - \mathbf{E}[f_n(\omega)]^2 \rightarrow 0$$

as $n \rightarrow \infty$, **iff** it is asymptotically uncorrelated with all **weighted majorities** $\text{Maj}_w(x_1, \dots, x_n) = \text{sign} \sum_{i=1}^n x_i w_i$. Also, not very slow decorrelation with all subset-majorities is enough for sensitivity.

Theorem. The left-right percolation crossing in an $n \times n$ square is noise sensitive, even with $\epsilon = \epsilon_n > c/\log n$.

Schramm & Steif 2005

Theorem. If $f : \Omega \rightarrow \mathbb{R}$ can be computed with a randomized algorithm with **revelment** δ (each bit is read only with probability $\leq \delta$), then

$$\sum_{S:|S|=k} \hat{f}(S)^2 \leq \delta k \|f\|_2^2.$$

For left-right crossing in $n \times n$ box on the hexagonal lattice, **exploration interface** with random starting point gives revelment $n^{-1/4+o(1)}$ (it has length $n^{7/4+o(1)}$, given by 2-arm exponent), while $\sum_{k \leq m} k \asymp m^2$, thus:

Theorem. Left-right crossing on the triangular lattice is noise sensitive under $\epsilon_n > n^{-a}$, with any $a < 1/8$. Even on square lattice, can take some positive $a > 0$.

The revelment is at least $n^{-1/2+o(1)}$ for *any* algorithm computing the crossing, hence this method can give only $n^{-1/4+o(1)}$ -sensitivity, far from the conjectured $\epsilon_n = n^{-3/4+o(1)}$.

The GPS approach, 2008

Although goal is to understand size, Gil Kalai suggested trying to understand entire distribution of \mathcal{S}_f . A strange random set of bits.

Effective sampling? If f is an effectively computable Boolean function, then there is an effective quantum algorithm for \mathcal{S}_f [Bernstein-Vazirani 1993].

For $\mathcal{S}_{Q,n}$ (left-right crossing in a conformal rectangle Q , mesh $1/n$), [Smirnov '01] + [Tsirelson '04] + [Schramm-Smirnov '11] implies that it has a conformally invariant scaling limit.

How to prove tightness for the size of strange random fractal-like sets?

Basic properties of the spectral sample

$$\text{For } A \subseteq V: \mathbf{E}[\chi_S \mid \mathcal{F}_A] = \begin{cases} \chi_S & S \subseteq A, \\ 0 & \text{otherwise.} \end{cases}$$

Therefore, $\mathbf{E}[f \mid \mathcal{F}_A] = \sum_{S \subseteq A} \hat{f}(S) \chi_S$, a nice projection.

Also, for $T \subseteq A$: $\mathbf{E}[f \chi_T \mid \mathcal{F}_{A^c}] = \sum_{S \subseteq A^c} \hat{f}(T \cup S) \chi_S$, hence

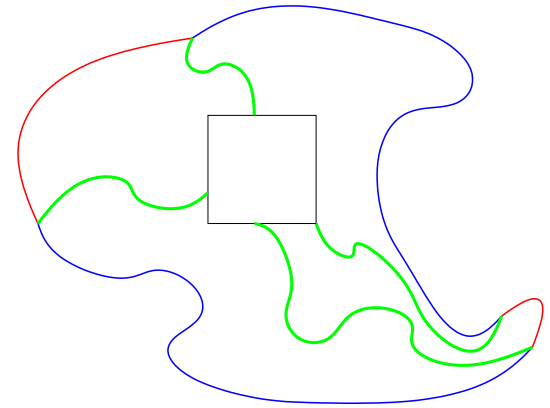
$$\mathbf{E}\left[\mathbf{E}[f \chi_T \mid \mathcal{F}_{A^c}]^2\right] = \sum_{S \subseteq A^c} \hat{f}(T \cup S)^2 = \mathbf{P}[\mathcal{S} \cap A = T].$$

This is the **Random Restriction Lemma** of **Linial-Mansour-Nisan '93**. E.g.,

$$\begin{aligned} \mathbf{P}[i, j \in \mathcal{S}_f] &= \mathbf{E}\left[\mathbf{E}[f \chi_{\{i,j\}} \mid \mathcal{F}_{\{i,j\}^c}]^2\right] \\ &= \frac{1}{4} \mathbf{P}[\omega|_{\{i,j\}^c} \text{ is such that } i, j \text{ each may be pivotal}] \\ &= \mathbf{P}[i, j \in \text{Piv}_f]. \end{aligned}$$

How does $[\mathcal{S}_n \cap B \mid \mathcal{S}_n \cap B \neq \emptyset]$ look like?

B as set has to be pivotal.



Strong Separation Lemma. For $d(B, \partial Q) > \text{diam}(B)$, conditioned on the 4 interfaces to reach ∂B , with *arbitrary starting points*, with a uniformly positive conditional probability the interfaces are well-separated around ∂B . Very bad separation is very unlikely. [Simple proof by [Damron-Sapozhnikov '09](#), following [Kesten '87](#). Also explained in Appendix to [GPS '11](#).]

Corollary 1. $\mathbf{P} \left[\mathcal{S}_n \cap B_r \neq \emptyset \right] \asymp \alpha_4(r, n)$.

Corollary 2. $\mathbf{E} \left[|\mathcal{S}_n \cap B_r| \mid \mathcal{S}_n \cap B_r \neq \emptyset \right] \asymp r^2 \alpha_4(1, r) \asymp \mathbf{E} |\mathcal{S}_r|$.

Self-similarity for left-right crossing of $n \times n$ square

$$\mathbf{E}|\mathcal{S}_n| = \mathbf{E}|\text{Piv}_n| \asymp n^2 \alpha_4(1, n) \stackrel{\Delta}{\asymp} n^{3/4+o(1)},$$

$$\mathbf{E}|\mathcal{S}_n(r)| := \mathbf{E}\left[\#\{r\text{-boxes } \mathcal{S}_n \cap B_r \neq \emptyset\}\right] \asymp \frac{n^2}{r^2} \alpha_4(r, n) \asymp \mathbf{E}|\mathcal{S}_{n/r}|,$$

$$\mathbf{E}\left[|\mathcal{S}_n \cap B_r| \mid \mathcal{S}_n \cap B_r \neq \emptyset\right] \asymp r^2 \alpha_4(1, r) \asymp \mathbf{E}|\mathcal{S}_r|.$$

Of course, $r^2 \alpha_4(1, r) \cdot \frac{n^2}{r^2} \alpha_4(r, n) \asymp n^2 \alpha_4(1, n)$, by **quasi-multiplicativity**.

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Similar to the **zero-set of simple random walk**: $\mathbf{E}|\mathcal{Z}_n| \asymp n n^{-1/2} = n^{1/2}$,

$$\mathbf{E}|\mathcal{Z}_n(r)| := \mathbf{E}\left[\#\{r\text{-intervals } \mathcal{Z}_n \cap I_r \neq \emptyset\}\right] \asymp \frac{n}{r} (n/r)^{-1/2} \asymp \mathbf{E}|\mathcal{Z}_{n/r}|,$$

$$\mathbf{E}\left[|\mathcal{Z}_n \cap I_r| \mid \mathcal{Z}_n \cap I_r \neq \emptyset\right] \asymp r r^{-1/2} \asymp \mathbf{E}|\mathcal{Z}_r|.$$

These results are related to the existence of scaling limits.

What concentration can we expect?

\mathcal{S}_n is very different from **uniform set** of similar density:
i.i.d. $\mathbf{P}[x \in \mathcal{U}_n] = n^{-5/4}$. Hence $\mathbf{E}|\mathcal{U}_n| = n^{3/4}$.

For large r ($\gg n^{5/8}$), this \mathcal{U}_n intersects every r -box;
for small r , if it intersects one, there is just one point there.

Concentration of size: roughly within $\sqrt{\mathbf{E}|\mathcal{U}_n|} = n^{3/8}$.

A bit more similar: for $i = 1, \dots, (n/r)^2$, i.i.d. $\mathbf{P}[X_i = r^{3/4}] = (n/r)^{-5/4}$,
 $X_i = 0$ otherwise. Then $S_{n,r} := \sum_i X_i$. Hence $\mathbf{E}|S_{n,r}| = n^{3/4}$.

For $r = n^\gamma$, size $|S_{n,r}|$ is concentrated within $n^{3/8(1+\gamma)}$, still $o(\mathbf{E}|S_{n,r}|)$.

For self-similar sets, we expect only **tightness around the mean**:
 $\mathbf{P}[0 < |\mathcal{S}_n| < \lambda \mathbf{E}|\mathcal{S}_n|] \rightarrow 0$ as $\lambda \rightarrow 0$, uniformly in n .

Proving tightness with a lot of independence

Assume we have the following ingredients, true for the zeroes:

$$(1) \quad \mathbf{P} \left[|\mathcal{Z}_n \cap I_r| > c \mathbf{E}|\mathcal{Z}_r| \mid \mathcal{Z}_n \cap I_r \neq \emptyset, \mathcal{F}_{[n] \setminus I_r} \right] \geq c > 0.$$

$$(2) \quad \mathbf{P} \left[|\mathcal{Z}_n(r)| = k \right] \leq g(k) \mathbf{P} \left[|\mathcal{Z}_n(r)| = 1 \right], \text{ with sub-exponential } g(k):$$

when the r -intervals intersected are scattered, have to pay k times to get to and leave them, and this cost is not balanced by combinatorial entropy.

Proving tightness with a lot of independence

Assume we have the following ingredients, true for the zeroes:

$$(1) \quad \mathbf{P} \left[|\mathcal{Z}_n \cap I_r| > c \mathbf{E}|\mathcal{Z}_r| \mid \mathcal{Z}_n \cap I_r \neq \emptyset, \mathcal{F}_{[n] \setminus I_r} \right] \geq c > 0.$$

$$(2) \quad \mathbf{P} \left[|\mathcal{Z}_n(r)| = k \right] \leq g(k) \mathbf{P} \left[|\mathcal{Z}_n(r)| = 1 \right], \text{ with sub-exponential } g(k):$$

when the r -intervals intersected are scattered, have to pay k times to get to and leave them, and this cost is not balanced by combinatorial entropy.

$$\mathbf{P} \left[0 < |\mathcal{Z}_n| < c \mathbf{E}|\mathcal{Z}_r| \right] = \sum_{k \geq 1} \mathbf{P} \left[0 < |\mathcal{Z}_n| < c \mathbf{E}|\mathcal{Z}_r|, |\mathcal{Z}_n(r)| = k \right]$$

$$\text{by (1):} \quad \leq \sum_{k \geq 1} (1 - c)^k \mathbf{P} \left[|\mathcal{Z}_n(r)| = k \right]$$

$$\text{by (2):} \quad \leq O(1) \mathbf{P} \left[|\mathcal{Z}_n(r)| = 1 \right] \asymp (n/r)^{1-3/2},$$

which, using $\lambda = \frac{c \mathbf{E}|\mathcal{Z}_r|}{\mathbf{E}|\mathcal{Z}_n|} \asymp (r/n)^{1/2}$, reads as $\mathbf{P} \left[0 < |\mathcal{Z}_n| < \lambda \mathbf{E}|\mathcal{Z}_n| \right] \asymp \lambda$.

But we know much less independence for \mathcal{S}_n

$$(1') \quad \mathbf{P} \left[|\mathcal{S}_n \cap B_r/3| > c \mathbf{E}|\mathcal{S}_r| \mid \mathcal{S}_n \cap B_r \neq \emptyset = \mathcal{S}_n \cap W \right] \geq c > 0,$$

for any W that is not too close to B_r .

Why only this negative conditioning? **Inclusion formula:**

$$\mathbf{P}[\mathcal{S}_f \subset U] = \sum_{S \subset U} \hat{f}(S)^2 = \mathbf{E} \left[\left(\sum_{S \subset U} \hat{f}(S) \chi_S \right)^2 \right] = \mathbf{E} \left[\mathbf{E}[f \mid \mathcal{F}_U]^2 \right].$$

From this, for disjoint subsets A and B ,

$$\begin{aligned} \mathbf{P}[\mathcal{S}_f \cap B \neq \emptyset = \mathcal{S}_f \cap A] &= \mathbf{P}[\mathcal{S}_f \subseteq A^c] - \mathbf{P}[\mathcal{S}_f \subseteq (A \cup B)^c] \\ &= \mathbf{E} \left[\mathbf{E}[f \mid \mathcal{F}_{A^c}]^2 - \mathbf{E}[f \mid \mathcal{F}_{(A \cup B)^c}]^2 \right] \\ &= \mathbf{E} \left[\left(\mathbf{E}[f \mid \mathcal{F}_{A^c}] - \mathbf{E}[f \mid \mathcal{F}_{(A \cup B)^c}] \right)^2 \right]. \end{aligned}$$

So, what are we going to do?

With quite a lot of work for both items,

$$(1') \quad \mathbf{P} \left[|\mathcal{S}_n \cap B_r/3| > c \mathbf{E}|\mathcal{S}_r| \mid \mathcal{S}_n \cap B_r \neq \emptyset = \mathcal{S}_n \cap W \right] \geq c > 0.$$

$$(2) \quad \mathbf{P} \left[|\mathcal{S}_n(r)| = k \right] \leq g(k) \mathbf{P} \left[|\mathcal{S}_n(r)| = 1 \right], \text{ with sub-exponential } g(k).$$

We could repeat (1') for many r -boxes only if “not enough points in one box” meant “we found nothing in that box”.

So, take an **independent random dilute sample**: $\mathbf{P}[x \in \mathcal{R}] = 1/\mathbf{E}|\mathcal{S}_r|$ i.i.d.

Then, $|\mathcal{S}_n \cap B_r/3|$ is small $\implies \mathcal{R} \cap \mathcal{S}_n \cap B_r/3 = \emptyset$ is likely,

and $|\mathcal{S}_n \cap B_r/3|$ is large $\implies \mathcal{R} \cap \mathcal{S}_n \cap B_r/3 \neq \emptyset$ is likely.

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and $|\mathcal{S}_n \cap B_r/3|$ is large $\implies \mathcal{R} \cap \mathcal{S}_n \cap B_r/3 \neq \emptyset$ is likely.

But $\mathbf{P} \left[\mathcal{S}_n \neq \emptyset = \mathcal{R} \cap \mathcal{S}_n \mid |\mathcal{S}_n(r)| = k \right]$ is still problematic conditioning.

A strange **large deviations lemma** solves the issue.

The strange large deviation lemma

Suppose $X_i, Y_i \in \{0, 1\}$, $i = 1, \dots, n$, and that $\forall J \subset [n]$ and $\forall i \in [n] \setminus J$

$$\mathbf{P}[Y_i = 1 \mid \forall_{j \in J} Y_j = 0] \geq c \mathbf{P}[X_i = 1 \mid \forall_{j \in J} Y_j = 0].$$

Then

$$\mathbf{P}[\forall_i Y_i = 0] \leq c^{-1} \mathbf{E}\left[\exp\left(-\frac{c}{e} \sum_i X_i\right)\right].$$

We use this with $X_j := 1_{\{\mathcal{S} \cap B_j \neq \emptyset\}}$ and $Y_j := 1_{\{\mathcal{S} \cap B_j \cap \mathcal{R} \neq \emptyset\}}$.

Proof: Instead of sequential scan, average everything together.

Choose $J \subset [n]$ randomly, Bernoulli($1-p$). Get $\mathbf{E}[Y p^Y] \geq c \mathbf{E}[X p^{Y+1}]$.

So, $\mathbf{E}[Z] \geq 0$, where $Z := (Y - cpX) p^Y$. Choose $p := e^{-1}$. Maximize Z over Y , and get the bound $Z \leq \exp(-1 - cX/e)$. Altogether, $ce^{-1} \mathbf{P}[Y = 0 < X] \leq \mathbf{E}[1_{X>0} \exp(-1 - cX/e)]$, and done.

Final result for the spectral sample

If $r \in [1, n]$, then $\{|\mathcal{S}_n| < \mathbf{E}|\mathcal{S}_r|\}$ is basically equivalent to being contained inside some $r \times r$ sub-square:

$$\mathbf{P}[0 < |\mathcal{S}_n| < \mathbf{E}|\mathcal{S}_r|] \asymp \alpha_4(r, n)^2 \left(\frac{n}{r}\right)^2.$$

In particular, on the triangular lattice Δ ,

$$\mathbf{P}[0 < |\mathcal{S}_n| < \lambda \mathbf{E}|\mathcal{S}_n|] \asymp \lambda^{2/3}.$$

The *scaling limit* of \mathcal{S}_n is a conformally invariant Cantor-set with Hausdorff-dimension $3/4$.

GPS (2010-12) proves that the **scaling limit of dynamical percolation** exists as a Markov process; for mesh $1/n$ the time-scale is $tn^{-3/4+o(1)}$. The above implies that this process is **ergodic**, with correlations decaying as $t^{-2/3}$.

Question 1: Can one build similar proofs for other Boolean functions?

Question 2: Self-similarity of Piv_n and \mathcal{S}_n is a lot of restriction on these random sets. And it's not only because of conformal invariance: **Gil Kalai** noticed that the spectral sample of **recursive 3-wise majority** is the leaves of a GW-tree! This phenomenon might be somehow general:

Influence-Entropy conjecture [Friedgut-Kalai 1996]: For some universal constant C , for any Boolean function f ,

$$\text{SpecEnt}(f) := \sum_{S \subset [n]} \hat{f}(S)^2 \log \frac{1}{\hat{f}(S)^2} \leq C \times$$
$$\times \text{Influence}(f) := \mathbf{E}|\mathcal{S}_f| = \mathbf{E}|\text{Piv}_f| = \sum_{S \subset [n]} \hat{f}(S)^2 |S|.$$

I.e., there is no log factor in the entropy as it would be in uniform.

I think I can do it for Piv_n , but not enough independence is known in \mathcal{S}_n .